### Algorithms Based On The Cauchy Point

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# The Cauchy Point

#### Perspective From Line Search

- Even when optimal step lengths are not used methods could be globally convergent.
- The step length  $\alpha_k$  needs to only satisfy fairly loose criteria.

#### Perspective For Trust-Region Method

- A similar imposition rather relation applies to trust-region methods as well.
- Even though the optimal solution to the sub-problem is seeked, it is enough to find an approximate solution *p<sub>k</sub>* within the trust region which gives <u>some sufficient reduction</u> to obtain <u>global convergence</u>.
- The sufficient reduction could be quantified in terms of Cauchy Point, which is denoted by p<sup>C</sup><sub>k</sub> and defined as follows:

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## Cauchy Point (Algorithm)

#### **Cauchy Point Calculation**

#### Step-I:

F

Find the vector  $p_k^S$  that solves a linear version of the trust region sub-problem i.e.

ind 
$$p_k^s$$
 s.t.  $||p_k^s|| \leq \Delta_k$  and  $p_k^s = rg \left\{ \min_{p \in \mathbb{R}^n} f_k + g_k^{\mathsf{T}} p 
ight\}$ 

#### Step-II:

Calculate the scalar  $\tau_k > 0$  that minimizes  $m_k(\tau p_k^s)$  subject to satisfying the trust-region bound i.e.

$$au_k = rg \left\{ \min_{ au \geq 0} m_k( au p_k^s) 
ight\} \qquad ext{s.t.} \ || au p_k^s|| \leq \Delta_k$$

Step-III:

Set 
$$p_k^C = \tau_k p_k^s$$
.

## Explicit Computation of The Cauchy Point

- Note that the problem in Step-I is a linear function.
- As a consequence p<sup>s</sup><sub>k</sub> chosen in the direction of the negative gradient should keep yielding reduction in the function value.
- That is to reduce the function one can move along the direction:



• As the function is linear the minimiser will lie at the boundary of the trust-region, giving

$$p_k^s = -rac{\Delta_k}{||g_k||}g_k.$$

• Now in Step-II we calculate  $\tau_k$ , we search for the minimiser of the model *m* in the direction of  $p_k^s$  (along the ray).

## Explicit Computation of The Cauchy Point

Towards this end we consider two cases

Case-I:  $g_k^T B_k g_k \leq 0$ Case-II:  $g_k^T B_k g_k > 0$ 

#### Case-I:

- The function  $m_k(\tau p_k^s)$  decreases monotonically with increasing  $\tau$  whenever  $g_k \neq 0$ .
- $f_k + \tau p_k^{s^T} g_k$  is decreasing as a consequence of the choice of  $p_k^{s^T}$ .
- Now since  $g_k^T B_k g_k \leq 0$  we have

$$f_k + \tau p_k^{s^T} g_k + \frac{1}{2} p_k^{s^T} B_k p_k^s$$

also decreases as  $p_k^{s^T} B_k p_k^s = \tau^2 \Delta_k^2 \frac{g_k^T B_k g_k}{||g_k||^2} \leq 0$ , when  $\tau$  increases.

## Explicit Computation of The Cauchy Point

• Therefore, the minimum is attained at simply the largest value that satisfies the trust-region bound for  $\tau_k$ , i.e.  $\tau_k = 1$ .

#### Case-II:

As  $g_k^T B_k g_k > 0$ ,  $m_k(\tau p_k^s)$  is a convex quadratic in  $\tau$ , so

•  $\tau_k$  is either the unconstrained minimiser of this quadratic i.e.

$$\tau_k = \frac{||g_k||^3}{\Delta_k g_k^\mathsf{T} B_k g_k}$$

• or, the boundary value 1. which ever comes first.

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#### Explicit Computation of The Cauchy Point

**Summary** 

$$p_k^C = -\tau_k \frac{\Delta_k}{||g_k||} g_k,$$

where

$$\tau_k = \begin{cases} 1, & \text{if } g_k^T B_k g_k \leq 0\\ \min\left(\frac{||g_k||^3}{\Delta_k g_k^T B_k g_k}, 1\right) & \text{otherwise.} \end{cases}$$

- The Cauchy step is inexpensive to calculate, no matrix factorisation are required.
- It is of crucial importance in deciding if an approximate solution of the trust-region sub-problem is acceptable.
- Specifically, a trust-region method will be globally convergent if its steps  $p_k$  give a reduction in the model  $m_k$  that is at least some fixed positive multiple of the decrease attained by Cauchy step.

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### The Cauchy Point

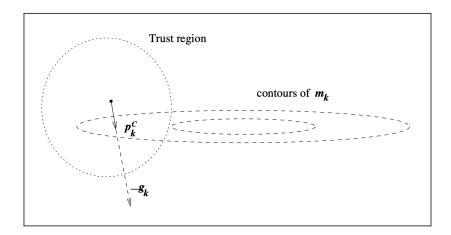


Figure: The Cauchy point for a subproblem in which  $B_k$  is positive definite. In this example,  $p_k^C$  lies strictly inside the trust region

# Improving On The Cauchy Point

- Why look any further, if  $p_k^C$  provides sufficient reduction for convergence and the cost of calculating it is so small.
- By always taking the Cauchy point as our step, we are simply implementing the steepest descent method with a particular choice of step length.
- Anyways, steepest descent performs poorly even for an optimal step length choice at each iteration.
- Similar issues one can raise as were done against steepest descent,

The Cauchy point doesn't depend very strongly on the matrix  $B_k$ , as it is used only for step length calculation.

## Improving On The Cauchy Point

- For rapid convergence  $B_k$  must be involved in the choice of direction as well. (if  $B_k$  contains valid information about the curvature of the function)
- Many trust-region algorithms compute the Cauchy point and then try to improve on it.
- The improvement strategy is often designed so that the full step length

$$p_k^B = -B_k^{-1}g_k$$

is chosen whenever  $B_k$  is positive definite and  $||p_k^B|| \leq \Delta_k$ .

When B<sub>k</sub> is the exact Hessian ∇<sup>2</sup>f(x<sub>k</sub>) or a quasi-Newton approximation, it can be expected to yield superlinear convergence.

- This method can be used when *B* is positive definite.
- Let us examine the effect of the (change in) trust-region radius Δ on the solution p\* (i.e. p\*(Δ)) of the trust-region sub-problem.
- When *B* is positive definite the unconstrained minimiser of the trust-region sub-problem is

$$p^B = -B^{-1}g$$

• If  $||p^B|| \leq \Delta$  (feasible) it is obviously the solution.

$$p^*(\Delta) = p^B, \qquad ||p^B|| \leq \Delta.$$

• When  $\Delta$  is very small relative to  $||p^B||$ , the restriction

 $||p|| \leq \Delta$ 

ensures that the quadratic term in m has little effect on the solution of the sub-problem.

 For such Δ, we can get an approximation to p(Δ) by simply omitting the quadratic term from m and writing

$$p^*(\Delta) pprox -\Delta rac{g}{||g||}, \qquad$$
 when  $\Delta$  is small.

 For intermediate values of Δ, the solution p\*(Δ) typically follows a curved trajectory like the one shown in the following figure.

For average values of  $\Delta$  the solution  $p^*(\Delta)$  usually follows a curved trajectory, as shown in the picture:

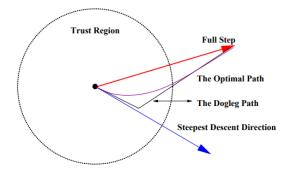


Figure: Exact trajectory and dogleg approximation.

- The dogleg method finds an approximate solution by replacing the curved trajectory for p<sup>\*</sup>(Δ) with a path consisting of two line segments.
- The first line segment runs from the origin to the minimiser of *m* along the steepest descent direction, which is

$$p^U = -\frac{g^T g}{g^T B g} g$$

- While the second line segment runs from  $p^U$  to  $p^B$ .
- Formally this trajectory is denoted by  $\tilde{p}(\tau)$  for  $\tau \in [0,2]$ , where

$$ilde{
ho}( au) = egin{cases} au p^U, & 0 \leq au \leq 1 \ p^U + ( au-1)(p^B-p^U), & 1 \leq au \leq 2. \end{cases}$$

Algorithms Based On The Cauchy Point The Dogleg Method

# The Dogleg Method

- The dogleg method chooses *p* to minimise the model *m* along this path, subject to the trust-region bound.
- The following lemma aids well towards the realisation of a minimum along the dogleg path (which can be easily found).

#### Lemma

Let B be positive definite. Then

- **1**  $||\tilde{p}(\tau)||$  is an increasing function of  $\tau$ , and
- 2  $m(\tilde{p}(\tau))$  is a decreasing function of  $\tau$ .

- It follows from the above lemma that the path  $\tilde{p}(\tau)$  intersects the trust-region boundary  $||p|| = \Delta$  at exactly one point if  $||p^{B}|| \geq \Delta$ , and nowhere otherwise.
- Since m is decreasing along the path, the chosen value of p will be at p<sup>B</sup> if ||p<sup>B</sup>|| ≤ Δ, otherwise at the point of intersection of the dogleg and the trust-region boundary.
- In the latter case, we compute the appropriate value of  $\tau$  by solving the following scalar quadratic equation:

$$|p^U + \tau * (p^B - p^U)|^2 = \Delta^2$$

$$||p^{U}||^{2} + 2\tau p^{U^{T}}(p^{B} - p^{U}) + \tau^{2}||p^{B} - p^{U}||^{2} = \Delta^{2}$$

Find the discriminant of the above equation. Find an explicit expression for  $\tau$  (Exercise).

# Summary of The Dogleg Method

- At each iteration, if the full step is within the trust region, it is used to update the current solution.
- If not, the algorithm searches for the minimum of the objective function along the steepest descent direction, known as Cauchy point.
- If the Cauchy point is outside of the trust region, it is truncated to the boundary of the latter and it is taken as the new solution.
- If the Cauchy point is inside the trust region, the new solution is taken at the intersection between the trust region boundary and the line joining the Cauchy point and the full step (dog leg step)

# Dogleg Point Algorithm

Compute 
$$p^{U}$$
  
and let  $x^{*} = x_{k} + p^{U}$   
if  $||p^{U}|| \ge \Delta_{k}$   
stop with  $x_{k+1} = x_{k} + \frac{\Delta_{k}}{||p^{U}||}p^{U}$   
else  
compute  $p^{B}$   
 $x^{*} = x_{k} + p^{B}$   
if  $||p^{B}|| \le \Delta_{k}$   
stop with  $x_{k+1} = x^{*}$ .  
else begin  
find  $\tau^{*}$  s.t.  $(|p^{U} + \tau * (p^{B} - p^{U})|^{2} = \Delta^{2})$   
stop with  $x_{k+1} = x_{k} + \tilde{p}(\tau^{*})$ 

# Dogleg Trivia

The name of the method derives from the resemblance between the construction of the dog leg step and the shape of a dogleg hole in golf.

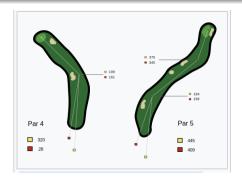


Figure: Typical doglegs. Left: "dogleg left". Right: "double dogleg".